

一、Mark each statement *true* or *false* (2 points each, 6 cents)

1. The same language token may be generated by many different regular expressions.
2. To any regular expression, we can find a context-free grammar defining the same language.
3. The LL(1) parsing algorithm parses an input string of tokens by tracing out the steps in a rightmost derivation.

二、Single Choice (1 points each , 5 cents)

1. The concept () is not related to the LL(1) parsing method.
[A] Left-factoring [B]. First set and follow set
[C.] Left recursion removal [D]. Shift and reduce
2. Which one below is not a part of a compiler? (这个题目不是太好, 就不扣分了)
[A] Symbol table [B] Assembler
[C] Code optimizer [D] Parser
3. In the production $A \rightarrow B \alpha C$, we have
[A] Follow (C) \subset Follow (A), First(B) \subset First(A)
[B] Follow (C) \subset Follow (A), First(A) \subset First(B)
[C] Follow (A) \subset Follow (C), First(B) \subset First(A)
[D] Follow (A) \subset Follow (C), First(A) \subset First(B)
4. IF one CFG grammar contains two non-terminals 'A','B' and two terminal 'a','b', where 'A' is the start symbol, then the Follow set of 'A' may be ()
[A] {a, b} [B] {a, b, \$} [C] {a, b, ϵ } [D] {a, b, B}
5. In the Top-Down Parsing, the action () will never be used.
[A] Shift [B] Match [C] Generate [D] Accept

三、question (39 cents)

1. Given the NFA for below for $0^*(01)^*0^*$, construct a minimum state DFA: (8 cents) (如果状态错的很多的话, 就给个基本分3分)
2. Given the follow grammar. (有的同学去掉了 M 再做下一步, 应该也算对的)
 $S \rightarrow L$
 $L \rightarrow MLb$
 $L \rightarrow a$
 $M \rightarrow \epsilon$
 (S is the start symbol.)
 Construct the LR(1) DFA for the grammar. (10 cents) (错一个状态扣 1 分)
3. (7 cents)

Give a RE and a CFG for:

$$L = \{x \in \{0,1\}^* \mid x \text{ starts and ends with different symbols} \}$$

4. Consider the following grammar of simplified C declarations:
 $\text{declaration} \rightarrow \text{type var-list}$
 $\text{type} \rightarrow \text{int} \mid \text{float}$
 $\text{var-list} \rightarrow \text{identifier, var-list} \mid \text{identifier}$
 (a) Left factor this grammar. (3 cents) (由于这个第一步没有做好影响下面的正确性, 后面扣分可以适当少一点的)
 (b) Construct First and Follow sets for the nonterminals of the resulting grammar.(6 cents) (错一个扣一分, 扣完为止, 注意: \$没有的不扣分)
 (c) Construct the LL(1) parsing table for the resulting grammar. (5 cents)
 (一个扣 0.5 分, 没有\$这一列的不扣分, 没有逗号的还是要扣分。)